# Maximum Entropy Approach to Interbank Lending: Towards a More Accurate Algorithm 

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#### Abstract

Banks loan money to each and borrow money from each other. To minimizing the risk caused by a possible default of one of the banks, a reasonable idea is to evenly spread the lending between different banks. A natural way to formalize this evenness requirement is to select the interbank amounts for which the entropy is the largest possible. The existing algorithms for solving the resulting constrained optimization problem provides only an approximate solution. In this paper, we propose a new algorithm that provides the exact solution to the maximum-entropy interbank lending problem.


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## 1 Formulation of the Problem

Interbank lending: formulation of the problem. Banks lend money to each other and borrow money from each other. To minimize the risk caused by a possible default of one of the banks, a reasonable idea is to spread the lending equally between as many banks as possible.

A natural way to describe such an even spread in precise terms is to make sure that the entropy of the resulting distribution is the largest possible; see, e.g., [1, 2, 3, 4, 5, 6, 7, 8, 9,
Interbank lending: formulation of the problem in precise terms.

- Let $a_{i}$ be the total amount of assets that the $i$-th bank is willing to lend.
- et $\ell_{j}$ be the amount of liabilities to cover which the $j$-th bank needs to borrow the money from other banks.

We assume that these amounts match, in the sense that we have a sufficient amount of available money to cover all the banks' liabilities:

$$
\begin{equation*}
S \stackrel{\text { def }}{=} \sum_{i=1}^{n} a_{i}=\sum_{j=1}^{n} \ell_{j}, \tag{1}
\end{equation*}
$$

where $n$ is the total number of banks.
For simplicity, we can divide all the amounts $a_{i}$ and $\ell_{j}$ by this common sum $S$, i.e., consider fractions of the overall lending amounts instead of the actual dollar values. After this division, the formula (1) takes a simplified form

$$
\begin{equation*}
\sum_{i=1}^{n} a_{i}=\sum_{j=1}^{n} \ell_{j}=1 \tag{1a}
\end{equation*}
$$

We want to describe, for every pairs $(i, j)$ of the banks $(i \neq j)$, the amount $x_{i j}$ that the $i$-th bank lends to the $j$-th one. We know the overall amount that each bank lends $a_{i}$, and the overall amount $\ell_{j}$ that each bank receives, thus:

$$
\begin{align*}
& \sum_{j=1}^{n} x_{i j}=a_{i} ;  \tag{2}\\
& \sum_{i=1}^{n} x_{i j}=\ell_{j} . \tag{3}
\end{align*}
$$

There are many different combinations of the values $x_{i j}$ that satisfy the constraints (2) and (3). Among these combinations, we would like to select a combination that maximizes the entropy

$$
\begin{equation*}
S=-\sum_{i, j} x_{i j} \cdot \ln \left(x_{i j}\right) . \tag{4}
\end{equation*}
$$

How this problem is solved now. At present, instead of directly solving the above optimization problem, researchers and practitioners use the following twostep approximate solution.

- First, they solve an auxiliary problem which is similar to the above but in which we are also allowing the values $x_{i i} \neq 0$. This auxiliary problem occurs in probability theory when:
- we know the marginal distributions $a_{i}$ and $\ell_{j}$, and
- we want to use the maximum entropy principle to select a joint distribution which is consistent with the given marginals.

A solution to this auxiliary problem is well known: it corresponds to the assumption that the corresponding random variables are independent:

$$
\begin{equation*}
x_{i j}^{(0)}=a_{i} \cdot \ell_{j} . \tag{5}
\end{equation*}
$$

- Then, we adjust these values to make sure that for the resulting values $x_{i j}$, we have $x_{i i}=0$. This is usually done by selecting the values $x_{i j}$ that maximize the conditional entropy

$$
\begin{equation*}
\sum_{i, j} x_{i j} \cdot \ln \left(\frac{x_{i j}}{x_{i j}^{(0)}}\right) . \tag{6}
\end{equation*}
$$

Limitations of the existing two-stage approach. While the above two-stage approach leads to a reasonable solution, it is only an approximate solution to the original maximum entropy problem. It is therefore desirable to come up with the exact (or at least more accurate) solution to this problem.

What we do in this paper. In this paper, we provide such a more accurate algorithm for the maximum-entropy interbank lending problem.

## 2 Analysis of the Problem and the Resulting Algorithm

Analysis of the problem. We want to maximize the entropy (4) under the constraints (2) and (3). By applying the Lagrange multiplier method, we can reduce this constrained optimization problem to the unconstrained problem of maximizing the following expression:

$$
\begin{equation*}
-\sum_{i j} x_{i j} \cdot \ln \left(x_{i j}\right)+\sum_{i=1}^{n} \lambda_{i} \cdot\left(\sum_{j=1}^{n} x_{i j}-a_{i}\right)+\sum_{j=1}^{n} \mu_{j} \cdot\left(\sum_{i=1}^{n} x_{i j}-\ell_{j}\right), \tag{7}
\end{equation*}
$$

for appropriate Lagrange multipliers $\lambda_{i}$ and $\mu_{j}$.

Differentiating the expression (7) with respect to $x_{i j}$ and equating the derivative to 0 , we conclude that

$$
-\ln \left(x_{i j}\right)-1+\lambda_{i}+\mu_{j}=0
$$

hence

$$
\ln \left(x_{i j}\right)=\left(\lambda_{i}-1\right)+\mu_{j}
$$

By applying $\exp (x)$ to both sides, we conclude that

$$
\begin{equation*}
x_{i j}=b_{i}^{\prime} \cdot c_{j}^{\prime} \tag{8}
\end{equation*}
$$

where we denoted $b_{i}^{\prime} \stackrel{\text { def }}{=} \exp \left(\lambda_{i}-1\right)$ and $c_{j}^{\prime} \stackrel{\text { def }}{=} \exp \left(\mu_{j}\right)$.
We can somewhat simplify this expression if we normalize the values $b_{i}^{\prime}$ by dividing them by the sum

$$
b^{\prime}=\sum_{j=1}^{n} b_{j}^{\prime}
$$

of these values, i.e., by considering the new values

$$
b_{i} \stackrel{\text { def }}{=} \frac{b_{i}^{\prime}}{b^{\prime}}
$$

for which

$$
\begin{equation*}
\sum_{i=1}^{n} b_{i}=1 \tag{9}
\end{equation*}
$$

Then, $b_{i}^{\prime}=b^{\prime} \cdot b_{i}$ and thus, the product $b_{i}^{\prime} \cdot c_{j}^{\prime}$ can be equivalently described as $b_{i} \cdot b^{\prime} \cdot c_{j}^{\prime}$, i.e., as

$$
\begin{equation*}
x_{i j}=b_{i} \cdot c_{j} \tag{10}
\end{equation*}
$$

for $i \neq j$, where we denoted $c_{j} \stackrel{\text { def }}{=} b^{\prime} \cdot c_{j}^{\prime}$.
For these values (10), the condition (4) takes the form

$$
\begin{equation*}
\ell_{j}=\sum_{i=1}^{n} x_{i j}=\sum_{i \neq j} b_{i} \cdot c_{j}=c_{j} \cdot \sum_{i \neq j} b_{i} \tag{11}
\end{equation*}
$$

Due to condition (1), the sum in the right-hand side of the formula (11) takes the form $1-b_{j}$, thus the condition (4) takes the form

$$
\begin{equation*}
c_{j} \cdot\left(1-b_{j}\right)=\ell_{j} \tag{12}
\end{equation*}
$$

Similarly, the condition (3) takes the form

$$
\begin{equation*}
a_{i}=\sum_{j=1}^{n} x_{i j}=\sum_{j \neq i} b_{i} \cdot c_{j}=b_{i} \cdot \sum_{j \neq i} c_{j} . \tag{13}
\end{equation*}
$$

Let us denote

$$
\begin{equation*}
C \stackrel{\text { def }}{=} \sum_{j=1}^{n} c_{j} . \tag{14}
\end{equation*}
$$

In this case, the sum in the right-hand side of the formula (13) takes the form $C-c_{i}$, hence the condition (13) takes the form

$$
\begin{equation*}
b_{i} \cdot\left(C-c_{i}\right)=a_{i} \tag{15}
\end{equation*}
$$

From (12), we can express $c_{i}$ as

$$
\begin{equation*}
c_{i}=\frac{\ell_{i}}{1-b_{i}} . \tag{16}
\end{equation*}
$$

substituting this expression into the formula (15), we conclude that

$$
\begin{equation*}
b_{i} \cdot\left(C-\frac{\ell_{i}}{1-b_{i}}\right)=a_{i} . \tag{16}
\end{equation*}
$$

Multiplying both sides of this equality by $1-b_{i}$, we conclude that

$$
\begin{equation*}
b_{i} \cdot C \cdot\left(1-b_{i}\right)-\ell_{i} \cdot b_{i}=a_{i} \cdot\left(1-b_{i}\right) \tag{17}
\end{equation*}
$$

Opening parentheses and moving all the term to the right-hand side, we get the following quadratic equation for determining $b_{i}$ :

$$
\begin{equation*}
b_{i}^{2}-b_{i} \cdot\left(C+a_{i}-\ell_{i}\right)+a_{i}=0 \tag{18}
\end{equation*}
$$

hence

$$
\begin{equation*}
b_{i}=\frac{C+a_{i}-\ell_{i}}{2} \pm \sqrt{\left(\frac{C+a_{i}-\ell_{i}}{2}\right)^{2}-a_{i}} \tag{19}
\end{equation*}
$$

For the large number of equal-size banks, the effect of $n$ terms $x_{i i}$ in the approximate solution is much smaller than the effect of $n^{2}-n$ other terns, so we should have $C \approx 1$ and $b_{i} \approx a_{i}$. In the formula (19), only the solution corresponding to minus has this asymptotics, so we should have minus:

$$
\begin{equation*}
b_{i}=\frac{C+a_{i}-\ell_{i}}{2}-\sqrt{\left(\frac{C+a_{i}-\ell_{i}}{2}\right)^{2}-a_{i}} \tag{20}
\end{equation*}
$$

Thus, once we know the value $C$ :

- we can find all the values $b_{i}$ by busing the formula (20),
- then we can find the values $c_{j}$ by using the formula (16), and
- hence, we can compute the values $x_{i j}=b_{i} \cdot c_{j}$.

We need to find the value $C$ for which, for the values $b_{i}$ from the formula (20), we have $\sum_{i=1}^{n} b_{i}=1$. Since $\sum a_{i}=\sum b_{j}$, the sum of the first (pre-square root) terms in the formula (20) is equal to $\frac{C \cdot n}{2}$, so the desired equality takes the form

$$
\begin{equation*}
\frac{C \cdot n}{2}-\sum_{i=1}^{n} \sqrt{\left(\frac{C+a_{i}-\ell_{i}}{2}\right)^{2}-a_{i}}=1 . \tag{21}
\end{equation*}
$$

Thus, we arrive at the following algorithm.
Resulting algorithm. First, we find the unknown value $C$ by solving the equation

$$
\begin{equation*}
\frac{C \cdot n}{2}-\sum_{i=1}^{n} \sqrt{\left(\frac{C+a_{i}-\ell_{i}}{2}\right)^{2}-a_{i}}=1 \tag{21}
\end{equation*}
$$

with a single unknown $C$.
Once we know $C$, we compute the auxiliary quantities

$$
\begin{equation*}
b_{i}=\frac{C+a_{i}-\ell_{i}}{2}-\sqrt{\left(\frac{C+a_{i}-\ell_{i}}{2}\right)^{2}-a_{i}} \tag{20}
\end{equation*}
$$

and

$$
\begin{equation*}
c_{i}=\frac{\ell_{i}}{1-b_{i}}, \tag{16}
\end{equation*}
$$

and compute the value $x_{i j}=b_{i} \cdot c_{j}$ for all pairs $i \neq j$.
These values are the exact maximum entropy solution to the interbank lending problem.

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